

Topological Strata of Weighted Complex Networks

Giovanni Petri^{1*}, Martina Scolamiero^{1,2}, Irene Donato^{1,3}, Francesco Vaccarino^{1,3}

1 ISI Foundation, Torino, Italy, **2** Dipartimento di Ingegneria Gestionale e della Produzione, Politecnico di Torino, Torino, Italy, **3** Dipartimento di Scienze Matematiche, Politecnico di Torino, Torino, Italy

Abstract

The statistical mechanical approach to complex networks is the dominant paradigm in describing natural and societal complex systems. The study of network properties, and their implications on dynamical processes, mostly focus on locally defined quantities of nodes and edges, such as node degrees, edge weights and –more recently– correlations between neighboring nodes. However, statistical methods quickly become cumbersome when dealing with many-body properties and do not capture the precise mesoscopic structure of complex networks. Here we introduce a novel method, based on persistent homology, to detect particular non-local structures, akin to *weighted holes* within the link-weight network fabric, which are invisible to existing methods. Their properties divide weighted networks in two broad classes: one is characterized by small hierarchically nested holes, while the second displays larger and longer living inhomogeneities. These classes cannot be reduced to known local or quasilocal network properties, because of the intrinsic non-locality of homological properties, and thus yield a new classification built on high order coordination patterns. Our results show that topology can provide novel insights relevant for many-body interactions in social and spatial networks. Moreover, this new method creates the first bridge between network theory and algebraic topology, which will allow to import the toolset of algebraic methods to complex systems.

Citation: Petri G, Scolamiero M, Donato I, Vaccarino F (2013) Topological Strata of Weighted Complex Networks. PLoS ONE 8(6): e66506. doi:10.1371/journal.pone.0066506

Editor: Renaud Lambiotte, University of Namur, Belgium

Received: January 26, 2013; **Accepted:** May 7, 2013; **Published:** June 21, 2013

Copyright: © 2013 Petri et al. This is an open-access article distributed under the terms of the Creative Commons Attribution License, which permits unrestricted use, distribution, and reproduction in any medium, provided the original author and source are credited.

Funding: GP is supported by the TOPDRIM project funded by the Future and Emerging Technologies program of the European Commission under Contract IST-318121. ID and MS are partly supported by Project Lagrange Ph.D. Grant. FV is partially supported by PRIN 2009 "Spazi di Moduli e Teoria di Lie". The funders had no role in study design, data collection and analysis, decision to publish, or preparation of the manuscript.

Competing Interests: The authors have declared that no competing interests exist.

* E-mail: giovanni.petri@isi.it

Introduction

Complex networks have become one of the prominent tools in the study of social, technological and biological systems [1–3]. In particular, weighted networks have been largely used to convey not only the presence but also the intensity of relations between nodes in a network. Real-world networks display however intricate patterns of redundant links with edge weights and node degrees usually ranging over various orders of magnitudes [4,5]. This makes very hard to extract the significant network structure from the background [6–9], especially in the case of very dense networks [10,11]. Alongside topological filtering methods [12,13], the typical approach to this problem is to choose a suitable threshold for the edge weights, e.g. global [10] or local [14], and study the reduced graph composed by only the edges of weight larger (smaller) than the threshold parameter. In any case, some properties of the original graph are inevitably lost under such transformation.

To avoid this pitfall, given a weighted network G we consider the set of all filtered networks, $\mathcal{F}(G)$, ordered by the descending thresholding weight parameter, in the spirit of *persistent homology* [15–18].

Persistent homology is a recent development in computational topology designed for robust shape recognition and data-discovery from high dimensional datasets [19]. It has found successful application in various fields, ranging from biological systems (e.g. brain correlation networks [20] and breast cancer diagnosis [15]), computer vision and sensor network coverage problems [15]

all the way to the analysis of large scale cosmological structure [22]. Its central device is the construction of a simplicial *filtration* of the original dataset: data points are usually embedded in a metric space in order to extract from their configuration a sequence of growing simplicial complexes, which approximates with increasing precision the original dataset. Studying the changes of the topological structure along such filtration provides a natural measure of robustness for the topological features emerging across different scales. In analogy to the metric example, we call the set $\mathcal{F}(G)$ *graph filtration*: considering the set of all filtered networks captures the link weights and connectivity structure over all weight scales, without the need to resort to any assumption on an eventual metric structure underlying the graph structure. The graph filtration of a network Ω is built following these steps :

- Rank the weights of links from ω_{max} to ω_{min} : the discrete parameter ϵ_t scans the sequence.
- At each step t of the decreasing edge ranking we consider the thresholded graph $G(\omega_{ij}, \epsilon_t)$, i.e. the subgraph of Ω with links of weight larger than ϵ_t .

Figure 1a provides a schematic illustration of the rank filtration. This approach preserves the complete topological and weight information, allowing us to focus on special mesoscopic structures: *weighted network holes*, that relate the network's weight-degree structure to its homological backbone.

A weighted network hole of weight ω is a loop composed by n nodes $i_0, i_1, i_2, \dots, i_{n-1}$, where all cyclic edges (i_l, i_{l+1}) (with $i_0 \equiv i_n$)

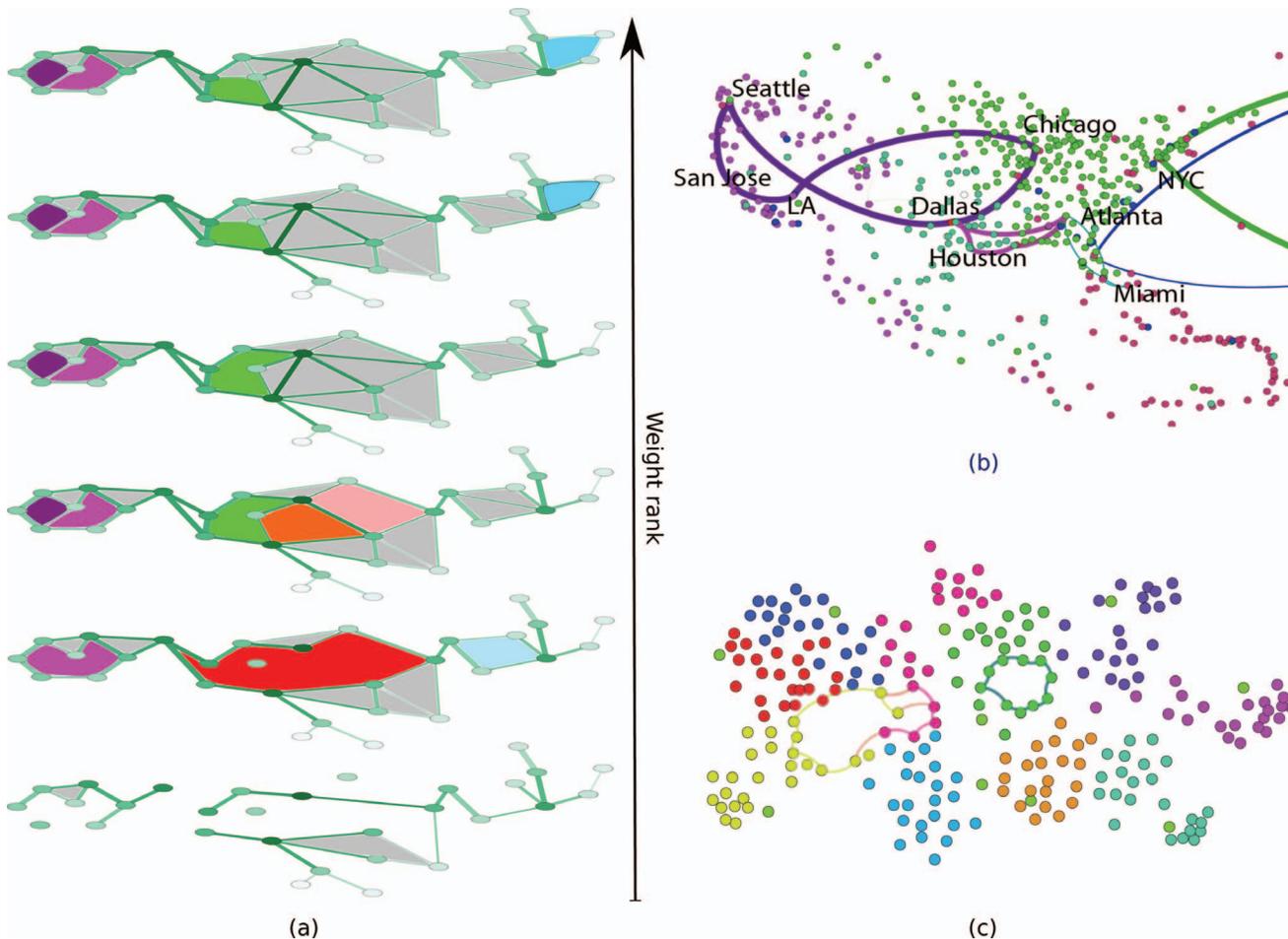


Figure 1. Weight rank clique filtration and homology of networks. (a) The weight rank filtration proceeds from the bottom up. Weighted holes (colored) and cliques (gray) appear as links are added. Weighted holes can branch into smaller holes, which have then independent evolution, persisting or dying along the filtration as links close them by 3-cliques. The cartoon shows two very long-persistence holes (violet and purple) appearing quite early and living until the end, while the largest hole (red) branches into three smaller holes, of only one survives to the end of the filtration (green). (b) A selection of weighted holes from the US air passenger network (year 2000). The node colors represent the best modularity partition of the entire network. The cycles are all long-persistence one, chosen to represent different behaviors: for example, the Chicago-Los Angeles-San Jose-Seattle cycle spans a large spatial distance, implying weaker connectivity across the cycle and within the region encompassed by the cycle, while the cycle going east from New York connects the east coast to three large European network and its persistence is due to the reduced connectivity due to the Atlantic Ocean. (c) A selection of the strongest cycles in the face-to-face contact network in a primary school (see SI for details on dataset). Node colors represent different classes in the school. Cycles are often found across communities, since by definition they probe the presence of holes among network regions. However, this is not the only information they convey. The cycle contained in a single community (green) testify the presence of peculiar contact geometries even within dense community structures.
doi:10.1371/journal.pone.0066506.g001

have weights $\geq \omega$, while all the other possible edges crossing the loop are strictly weaker than ω . We focus on this special class of subgraphs, because formally such weighted holes represent the generators of the first homology group, H_1 , of the clique complex of the graph thresholded by weight ω (see Materials and Methods). The aim of this paper is to characterize the evolution of these generators along the network filtration. As we swipe the network from the largest to the smallest weights, network holes appear and potentially close.

By unearthing their properties, we obtain the main contribution of this paper: the statistical features of weighted network holes yield a classification of real-world networks in two classes, depending on the compatibility or lack thereof with null models generated by graph randomisations. Furthermore, this classification is defined by mesoscopic homological structures that cannot be reconduced to local properties alone.

The method used for the classification itself, which we call *weighted clique rank homology*, is the second novel main contribution of this paper. It allows to recover complete and accurate long-range information from noisy redundant network data, by building on persistent homology [16], a recent theory developed in computational topology [17], which we extend to the case of networks.

Each weighted hole g is characterized by three quantities: its birth index β_g , its persistence p_g and its length λ_g . After ranking links in a descending order according to their weights, the birth index of a hole is the rank t of its weight ω . As we proceed adding links to the filtration in ranking order, it is possible that a link with rank $t' > t$ will appear and cross the hole. We call this closure of the weighted hole, or *death* δ_g . The persistence p_g is the interval between the birth and death of g , $p_g = \delta_g - \beta_g = t' - t$. Finally, the length λ_g is the number of links composing g .

Similarly to stratigraphy, each step of the filtration is a topological stratum of the network, where the edge weight rank plays the role of depth. Intuitively, g can then be thought as an underground cavity, hidden in the link-weight fabric of the network, and β_g , p_g and λ_g as its maximal depth, vertical size and girth respectively.

Results

Homological Network Classes

We applied this analysis to various social, infrastructural and biological networks (see SI for a detailed list). In order to compare datasets, indices are normalized by the corresponding filtration length (maximal rank) T , so that all β_g , δ_g , and thus p_g , vary in the unit interval. In addition, we compared each dataset with two randomized versions, obtained by weight reshuffling and edge-swapping respectively. While both randomisations preserve the weight and degree sequences (and the relative distributions ($p(k)$ and $p(\omega)$)), the first one redistributes only the edge weights and is meant to destroy weight correlations, preserving the joint degree distribution $p(k, k')$ and thus the degree assortativity. The second instead randomizes the network through double-edge swaps, preserving $p(k)$ and $p(\omega)$ but destroying both weight and degree correlations [23]. We stress that, as the degree and weight sequences are preserved in the randomisations, they cannot account for the differences in the observed homology.

The statistical distributions obtained for the $\{\beta_g\}$, $\{p_g\}$ and $\{\lambda_g\}$ for H_1 cycles highlight a natural division of the analysed networks in two broad classes (Fig. 2):

Class I networks. cycle distributions are markedly different from the randomized versions (cycles display shorter persistence times, earlier and broader birth distributions and very short lengths as compared to their randomized versions);

Class II networks. cycle distributions are very close to their random versions (late appearance, short persistences, long cycles).

The short cycles of Class I networks nest hierarchically and appear and die over all scales while those in the randomized counterparts are born uniformly along the filtration but are more persistent, producing largely hollow network instances. The implications are twofold. Since cycles represent weaker connectivity regions, this results in class I networks being more *solid* than the randomized versions, while class II networks resemble more closely the randomized instances. Second, since the cycle abundance ratio between real and random instances is the same in the two groups, the differences between class I and II does not depend on cycle abundance, but rather on their properties.

This can be seen easily by compressing the whole information within two scalar metrics which do not depend on the number of generators in a given network filtration. We define the *network hollowiness* h_i and the *chain-length normalized hollowiness* \tilde{h}_i as:

$$h_k = \frac{1}{N_{g_k}} \sum_{g_k} \frac{p_{g_k}}{T} \quad (1)$$

$$\tilde{h}_k = \frac{1}{N_{g_k}} \sum_{g_k} \frac{\lambda_{g_k} p_{g_k}}{N} \frac{p_{g_k}}{T} \quad (2)$$

where $\{g_k\}$ is the set of generators of the k -th homological group H_k and $N_{g_k} = \dim H_k$ their number. The first is a measure of the average generator persistence, while the second weights generators according to both their length and persistence. Table 1 reports the

values for h_1 and \tilde{h}_1 . Class I networks have lower hollowiness values as compared to their randomized versions, while class II ones show comparable values.

Interestingly, the hollowiness values for the H_2 generators mostly vanish for the randomized instances (Table 1), as opposed to the case of real networks. It appears that, while persistent one-dimensional cycles are more easily generated in the randomized instances, higher forms of network coordination, e.g. H_2 generators (akin to two-dimensional surfaces bounding three-dimensional voids), do not only display different properties in comparison to the real network, but are instead wiped away. These findings hint therefore to the presence of higher order coordination mechanisms in real world networks.

Naturally, the two network classes do not represent a binary taxonomy and should be considered as two extremes of a range over which networks are distributed. For example, we find networks that interpolate between these classes, e.g. the online messages network has short persistence intervals, but also late cycle appearances and short length cycles. However, classes do not appear to display uniform behavior for local and two-body quantities: degree- and weight-distributions and correlations are mixed within the same group and do not provide a direct answer for the nature of the two classes. Similarly, a recently proposed measure of structural organisation, *integrativeness* [24], which measures the neighborhood overlap around strong links, does not provide insights to explain class I, since within the latter one finds both integrative and dispersive networks.

Finally, the classes do not show a consistent pattern in *assortativity*: for example, class I includes the gene network (assortative) and the airport networks (disassortative), while class II includes the assortative co-authorship networks and the disassortative Twitter data. Therefore, assortativity cannot be the discriminating factor between classes.

Higher Order Organization

Because homology is essentially a non-local property, it was expectable that the local measures mentioned would not be able to explain the observed homological patterns. Network homology can be seen in fact as the weighted complement to the *perturbative dK-series* approach [8]: the latter proceeds by successive bottom-up constraints on k -body correlations, rapidly becoming very cumbersome, while our method returns the complete superposition of the network's degree and weight correlation layers in a non-perturbative (top-down) fashion.

A simple artificial network helps illustrating this point: Random Geometric Graphs (RGG) have been recently shown to display long-range many-body correlations [25,26]. We find also that they have homological structures reminding of class I networks (Fig. 2a, b and c) and the same relation to their randomized versions. Class I networks are the result of high-order coordination in a similar way. This is supported also by the presence in real networks and RGGs of higher homology generators, which require elaborate coordination patterns in order to appear. While these cycles almost disappear in randomized versions of real-world networks, they are present in the case of RGGs.

For the latter and the airports, this organisation can be thought as the result of the non-local constraint imposed by the metric of the underlying space [27]. Although spatial constraints are harder to fathom for social and genetic systems, alternative explanations are possible: for example, the homological structure of the observed online communication and gene networks can be thought as stemming from group interactions among people (e.g. mailing lists, multi-user mails) and biological functions (e.g.

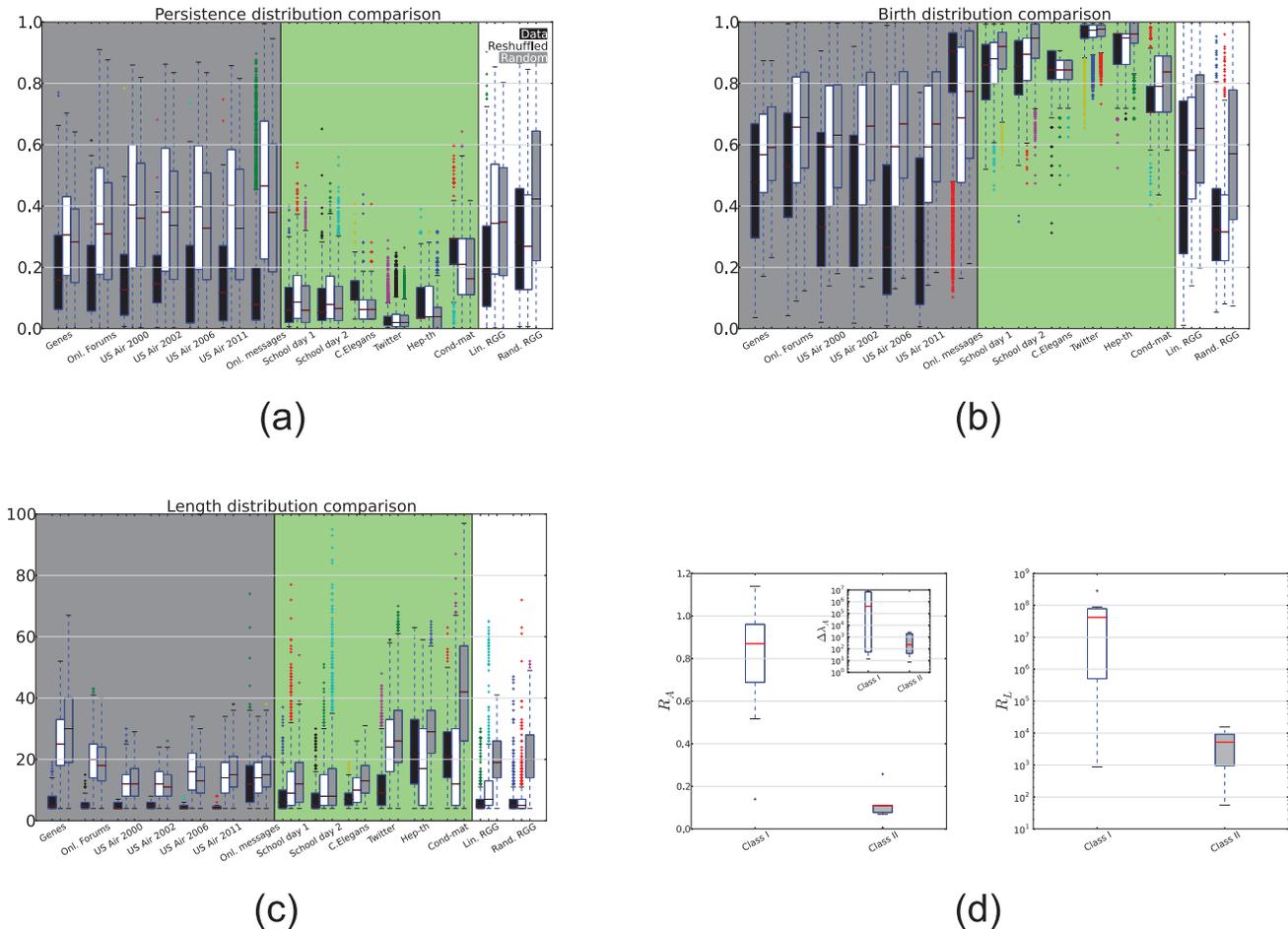


Figure 2. Statistical and spectral properties of H_1 generators. Box plots of the distributions of persistences $\{p_g\}$ (panel a), births $\{\beta_g\}$ (panel b) and lengths $\{\lambda_g\}$ (panel c) for the 1d cycles (H_1 generators) of real networks (black), reshuffled (white) and randomized (gray). The gray and green shaded areas identify the two network classes described in the main text: class I is significantly different from the random expectations, with shorter, less persistent cycles that appear across the entire filtration; class II networks are not significantly different from the random versions, with long cycles and late birth times in the filtration. The characteristics of class I networks imply a stratification of cycles that betrays the presence of large, non-local organisation in the network structure, which is not present in class II networks. For comparison, an example of RGG network (600 nodes in the unitary disk, linking distance 0.01), known to have higher order degree correlations, had edge weights set according to $\omega_{ij} \propto (k_i k_j)^\theta$, with $\theta = 1$ (linearly correlated weight RGG) and $\theta = 0$ (random weight RGG). In both cases, the distributions of cycles' properties resemble closely those of class I networks. Panel d finally reports the distribution of adjacency spectral gaps $\Delta\lambda_A$ and R_A (left plot) and the Laplacian eigenratio R_L (right plot). All the quantities show significant ($p < 0.05$) differences between the two classes, implying that the homological structure affect the dynamical properties of networks, e.g. the synchronizability threshold. doi:10.1371/journal.pone.0066506.g002

pathways) respectively, which provide an underlying non-local mechanism for the emergence of homological patterns.

Further evidence of this behavior can be found by zooming on specific cycles which convey information about underlying constraints hidden in the network weight-link connectivity patterns. For example, the cycle structure of the air passenger network detects the expected reduced connectivity over oceans in the form of strong persistent cycles—and the strong backbone of US airport hubs, which is then filled by the local (intra-community) links (Fig. 1b). Another example can be found in the school children's face-to-face contact network. As expected we find the most significant cycles to link together different school classes (yellow and pink cycles in Fig. 1c). However, we also find that a school class (green nodes), despite being both a network community and 3-clique component [28], is characterized by a strong internal H_1 generator, which might be reflecting peculiar social dynamics

coming from same-gender biases, different seating arrangements or schedules for part of the class [29].

Spectral Correlates of Homology Classes

At the opposite extreme of local quantities lie the spectral properties of networks. It is very important therefore to investigate whether it is possible to highlight peculiar spectral signatures of the two classes. Network eigenvalues, especially those of the Laplacian matrix, figure prominently in a number of applications, ranging from spectral clustering [30] to the propensity to synchronize of a set of oscillators distributed on the nodes [31]. Given a graph G , we denote its adjacency matrix $A(G)$ and its Laplacian matrix as $L(G) = D - A(G)$, where $d_{ij} = \delta_{ij} \sum_k a_{ik}$. For a symmetric network with N nodes, $A(G)$ has a set of real eigenvalues $\lambda_1 \geq \lambda_2 \geq \dots \lambda_{N-1} \geq \lambda_N$. The spectral gap $\Delta\lambda_A = \lambda_1 - \lambda_2$, and its normalized version, $R_A = \frac{\lambda_1 - \lambda_2}{\lambda_2 - \lambda_N}$, effectively measure how far the

leading eigenvalue lies in comparison to the bulk of the eigenvalue distribution [32].

Interestingly, we find that class I networks have significantly larger spectral gaps ($p < 0.05$ comparing the distributions) than class II networks (in Fig. 2d and Table 2 for information on individual datasets). Despite being somewhat neglected in the complex networks literature, $\Delta\lambda_A$ has been linked to the notion of natural connectivity [33]: it encodes spectral information about network redundancy in terms of the number of closed paths and is defined as $\bar{\lambda} = \log \left[\frac{1}{N} \sum_{i=1}^N e^{\lambda_i} \right]$. Rewriting $\bar{\lambda} = \lambda_1 + \log \left[\frac{1}{N} (1 + \sum_{i=2}^N e^{\lambda_i - \lambda_1}) \right]$, it is easy to see that for large gaps all the terms in the sum are exponentially suppressed and therefore $\bar{\lambda}$ is essentially dominated by the leading adjacency eigenvalue modulo a size effect, $\bar{\lambda} \sim \lambda_1 - \log N$. This result is consistent with the nested cycle structure that we highlighted in class I. More importantly, we find a difference between the two classes in the topological constraints to synchronization processes. For the Laplacian $L(G)$, label the set of eigenvalues $0 = \lambda_1 < \lambda_2^L \leq \lambda_3^L \leq \dots \leq \lambda_N^L$ and define the Laplacian eigenratio $R_L = \frac{\lambda_N^L}{\lambda_2^L}$. Barahona and Pecora [34] showed that a set of dynamical systems, placed on the network's nodes and coupled according to the graph adjacency with a global coupling σ , has a linearly stable synchronous state if

$$R_L < \beta \quad (3)$$

where β is a purely dynamical parameter. This inequality implies that networks displaying very large R_L are hard (or impossible) to synchronize. Panel IVb of Fig. 2 shows again a significant difference between the two classes: class I networks have much larger eigenratios, making them hardly synchronizable.

Our results show therefore a deep connection between the homological network structure, the network spectral properties and their implications on network dynamics. Indeed, the role of mesoscopic structures in the stability and evolution of dynamical systems on networks is gradually emerging, as shown for example by recent work based on the concepts of basic symmetric subgraphs and their legacy eigenvalues in the global network spectrum [35], and is indeed being shaped by algebraic methods, well suited to capture the geometric information hidden within the network fabric.

Conclusions

Hitherto, the homological structure of weighted networks could not be systematically studied. Our method, grounded in computational topology, allows to probe multiple layers of organized structure. It highlighted two classes of network distinguished by their homological features, which we interpreted as caused by differences in the higher order networks organisations that are not captured by (quasi)local approaches.

Among the many possible applications, two very relevant ones for social and infrastructural networks are the study of the weighted rich club's geometry beyond the aggregate measure [23,36], and the generalisation of network embedding models to include homological information [37]. Furthermore, the two classes displayed also a marked difference in their spectral gap distributions and in particular in the values of the algebraic connectivity, implying that the different homological structures are correlated with different synchronizability thresholds.

This work therefore provides a stepping stone towards understanding the coupling between network dynamical processes and the network's homology.

Finally, the filtration's construction rule is flexible and can be readily adapted to other problems. Similarly to changing goggles, different edge metrics can be used (e.g. betweenness or salience [38]), the thresholding method varied (e.g. local thresholding [14]) or the filtration promoted to a filtering on two quantities (e.g. edge weight and time in a temporal network) using *multi-persistent* homology [39].

Materials and Methods

Datasets

The dataset analysed in this paper cover a broad range of fields, spanning social, infrastructural and biological networks. Figures S1–S15 in the File S1 report the analysis for the individual datasets as opposed to the class-aggregate of figure 2.

In detail, they are:

US air passenger networks. The networks refer to the years 2000, 2002, 2006 and 2011. The years were chosen to provide snapshots of the air traffic situation at 4–5 years intervals, plus one extra (year 2000) just before the events of 9/11 which significantly affected the air transportation industry. The data used are publicly available from the website of the Bureau of Transportation Statistics (<http://www.transtats.bts.gov/>). Individual flights between airports were aggregated on routes as defined by origin and destination cities. The weight reported is the yearly aggregated passenger traffic.

C.Elegans. The network is available at <http://cdg.columbia.edu/cdg/datasets> and reports a weighted, directed representation of the C. Elegans's neuronal network [40]. The network was symmetrized by summing the weights present on edges between the same nodes (given ω_{ij} and ω_{ji} , $\omega_{ij}^{symm} = \omega_{ji}^{symm} = \omega_{ij} + \omega_{ji}$).

Online messages and forums. The online messages network consists of messages in a student online community at University of California [41]. The online forum network refers to the same online community, but focuses on the activity of users in public forums, rather than on private messages [42]. Both networks are publicly available online at Tore Opsahl's website (<http://toreopsahl.com/datasets/>).

Gene network. The gene interaction network used in the paper is a sampling of the complete human genome dataset available from the University of Florida Sparse Matrix Collection. Each node is an individual gene, while the edges correlates the expression level of a gene with that of the genes (using a NIR score [43]). The node set of the analysed network was obtained by randomly choosing an origin node, then adding its neighborhood to the node set; the neighborhoods of the newly added nodes were then added to the node set recursively until a given number of nodes was obtained (in the case used the target number of nodes was $N = 1300$). Then all the edges present in the original network between the nodes in the node set were added, effectively taking a connected subgraph of the original network. To reduce the computational complexity due to the large density of the graph, the weighted clique filtration was stopped at an edge weight of 0.09 (similarly to the choice made in [24]).

Twitter. The dataset consists of a network of mentions and retweet between Twitter users and is available online on the Gephi dataset page (<http://wiki.gephi.org/index.php/Datasets>). Weights are proportional to the number of interactions between a pair of users.

School face-to-face contact network. The dataset contains two days of recorded face-to-face interactions in a primary school.

Table 1. Summary of hollowness values.

| Dataset (class) | h_1 | \tilde{h}_1 | h_1^{sh} | \tilde{h}_1^{sh} | h_1^{rnd} | \tilde{h}_1^{rnd} | h_2 | \tilde{h}_2 |
|--------------------|--------|---------------|---------------------|----------------------|---------------------|-----------------------|---------|---------------|
| Genes(I) | 0.515 | 0.003 | 0.020 ± 0.001 | 0.0007 ± 0.00001 | 0.0151 ± 0.0004 | 0.00023 ± 0.00005 | 0.35 | 0.006 |
| Online forums(I) | 0.175 | 0.001 | 0.355 ± 0.005 | 0.007 ± 0.001 | 0.325 ± 0.005 | 0.007 ± 0.001 | 0.02 | 0.0003 |
| US Air 2000(I) | 0.160 | 0.001 | 0.405 ± 0.005 | 0.0065 ± 0.0007 | 0.358 ± 0.006 | 0.0060 ± 0.0005 | 0.02 | 0.0003 |
| US Air 2002(I) | 0.186 | 0.0008 | 0.39 ± 0.01 | 0.0037 ± 0.0003 | 0.34 ± 0.01 | 0.0034 ± 0.0003 | 0.23 | 0.002 |
| US Air 2006 (I) | 0.167 | 0.0005 | 0.398 ± 0.005 | 0.0036 ± 0.0005 | 0.348 ± 0.008 | 0.0032 ± 0.0003 | 0.165 | 0.001 |
| US Air 20011(I) | 0.181 | 0.0006 | 0.41 ± 0.01 | 0.0034 ± 0.0002 | 0.35 ± 0.01 | 0.0033 ± 0.0003 | 0.076 | 0.0007 |
| Online messages(I) | 0.21 | 0.0014 | 0.190 ± 0.002 | 0.0017 ± 0.0001 | 0.185 ± 0.002 | 0.0015 ± 0.0001 | 0.02 | 0.0003 |
| School day 1 (II) | 0.088 | 0.0034 | 0.113 ± 0.002 | 0.007 ± 0.001 | 0.093 ± 0.002 | 0.006 ± 0.001 | 0.015 | 0.0012 |
| School day 2 (II) | 0.090 | 0.0033 | 0.115 ± 0.002 | 0.0065 ± 0.0005 | 0.098 ± 0.003 | 0.0089 ± 0.0008 | 0.01412 | 0.00095 |
| C. elegans (II) | 0.0784 | 0.002 | 0.0745 ± 0.0017 | 0.001 ± 0.0001 | 0.0896 ± 0.0023 | 0.0041 ± 0.0005 | 0.058 | 0.002 |
| Twitter (II) | 0.03 | 0.0001 | 0.030 ± 0.001 | 0.0002 ± 0.0001 | 0.029 ± 0.001 | 0.0002 ± 0.0001 | 0.01 | 0.0001 |
| Hep-th (II) | 0.08 | 0.0002 | 0.075 ± 0.001 | 0.0002 ± 0.0001 | 0.0508 ± 0.0003 | 0.0002 ± 0.0001 | – | – |
| Cond-mat (II) | 0.26 | 0.0004 | 0.20 ± 0.003 | 0.0002 ± 0.0001 | 0.180 ± 0.002 | 0.0005 ± 0.0001 | – | – |
| Lin. RGG | 0.227 | 0.003 | 0.368 ± 0.005 | 0.006 ± 0.001 | 0.355 ± 0.002 | 0.012 ± 0.001 | 0.28 | 0.006 |
| Ran. RGG | 0.3 | 0.0041 | 0.299 ± 0.005 | 0.0045 ± 0.0002 | 0.649 ± 0.40 | 0.015 ± 0.001 | 0.115 | 0.003 |

Summary of hollowness values. For each dataset, we report the values of the *hollowness* h_1 and *cycle-length normalized hollowness* \tilde{h}_1 for H_1 cycles for real networks and their randomisations (*sh* and *rnd*). Most networks (class I in particular) show lower values than for their randomized versions. We also report the values of the *hollowness* h_2 and *cycle-length normalized hollowness* \tilde{h}_2 for H_2 cycles for real networks. The values for the randomized networks are not reported as –strikingly– the randomisations do not inline any higher homology, while almost all real networks inline positive values of the H_2 hollowness. doi:10.1371/journal.pone.0066506.t001

Each node represents a child, with the edge weight between two nodes being proportional to the amount of time the two children spent face to face. We analysed the two days separately, yielding two networks. The dataset has been collected by the Sociopattern project (<http://www.sociopatterns.org/>) and analysed in [29].

Co-authorship networks. The networks analysed are the weighted co-authorship networks of the Condensed Matter E-print Archive between 1995 and 1999 (cond-mat) and the High-Energy Theory E-print Archive between 1995 and 1999 (hep-th) [44].

The graph edgelist used in the paper are available online as part of the code package we developed [45].

Finally, for comparison we use Random Geometric Graphs (RGG) [46,47], which are simple models of spatial networks: a RGG is generated by sprinkling N of nodes randomly on a metric space that acts as a substrate (usually a disk of unitary radius or a square with identified edges), and then linking nodes that are closer than a given linking distance d .

The networks analysed in this article are undirected and weighted, because the weighted clique filtration finds a natural application in such case. However, schemes for directed networks can be easily devised and tailored to specific case studies, e.g. one could adopt the definition used in the directed clique percolation method [48] in order to associate network structures to simplices.

Persistent Homology

The method we use to uncover weighted holes is persistent homology of the weight clique rank filtration. In this section we will briefly explain persistent homology and its realization through the weight rank clique filtration.

Persistent homology is a technique from computational algebraic topology that can be viewed as parametrized version of simplicial homology [49]. The two definitions needed for simplicial homology are those of *simplicial complex* and *homology*. A

simplicial complex is a non empty family X of finite subsets, called faces, of a vertex set with the two constraints:

- a subset of a face in X is a face in X ,
- the intersection of any two faces in X is either a face of both or empty.

We assume that the vertex set is finite and totally ordered. A face of $n + 1$ vertices is called n –face and denoted by $[p_0, \dots, p_n]$. The interpretation of low dimensional faces is intuitive: a 0–face is a vertex, a 1–face is a segment, a 2–face is a full triangle, a 3–face is a full tetrahedron. The dimension of a simplicial complex is the highest dimension of the faces in the complex.

Morphism between simplicial complexes are called simplicial maps. A simplicial map is a map between simplicial complexes with the property that the image of a vertex is a vertex and the image of a n –face is face of dimension $\leq n$.

Simplicial Homology with coefficients in a field is a functor from the category of simplicial complexes to the category of vector spaces [49]. Homology of dimension n assigns to each simplicial complex X , the vector space $H_n(X)$ of n -cycles modulo boundaries and to every simplicial map $X \xrightarrow{f} Y$ the linear map $H_n(f) : H_n(X) \rightarrow H_n(Y)$.

The construction that leads to the vector space H_n is the following. Given a simplicial complex X of dimension d , consider the vector spaces C_n on the set of n –faces in X for $0 \leq n \leq d$. Elements in C_n are called n –chains. The linear maps sending a n –face to the alternate sum of its $(n - 1)$ –faces

$$\partial_n : C_n \rightarrow C_{n-1}$$

Table 2. Summary of spectral quantities values.

| Dataset (class) | R_A | $A\lambda_A$ | R_L |
|--------------------|--------|------------------|------------------|
| Genes(I) | 1.14 | 14.6 | 873 |
| Online forums(I) | 0.5 | $4 \cdot 10^5$ | $3.4 \cdot 10^5$ |
| US Air 2000(I) | 0.868 | $6.9 \cdot 10^6$ | $6.7 \cdot 10^6$ |
| US Air 2002(I) | 0.872 | $6.3 \cdot 10^6$ | $2.8 \cdot 10^7$ |
| US Air 2006 (I) | 0.958 | $7.7 \cdot 10^6$ | $4.2 \cdot 10^7$ |
| US Air 20011(I) | 0.941 | $6.9 \cdot 10^6$ | $8.9 \cdot 10^7$ |
| Online messages(I) | 0.14 | $1.1 \cdot 10^4$ | $6.7 \cdot 10^4$ |
| School day 1 (II) | 0.11 | $2.5 \cdot 10^3$ | 56 |
| School day 2 (II) | 0.08 | $2.3 \cdot 10^3$ | 110 |
| C. elegans (II) | 0.25 | 76 | $1.8 \cdot 10^3$ |
| Twitter (II) | 0.11 | 370 | $1.5 \cdot 10^4$ |
| Hep-th (II) | 0.11 | 7.4 | $9.6 \cdot 10^3$ |
| Cond-mat (II) | 0.005 | 0.24 | $5.2 \cdot 10^3$ |
| Lin. RGG | 0.0034 | 34 | 836 |
| Ran. RGG | 0.018 | 54 | 255 |

Summary of spectral quantities. For each dataset, we report the values of R_A , $A\lambda_A$ and R_L . The two classes inline different spectral properties, with particular reference to R_A which is related to the network expansion property. doi:10.1371/journal.pone.0066506.t002

$$[p_0, \dots, p_n] \rightarrow \sum_{i=0}^n (-1)^i [p_0, \dots, p_{i-1}, p_{i+1}, \dots, p_n].$$

shares the property $\partial_{n-1} \circ \partial_n = 0$.

The subspace $\ker \partial_n$ of C_n is called the vector space of n -cycles and denoted by Z_n . The subspace $\text{Im} \partial_{n+1}$ of C_n , is called the vector space of n -boundaries and denoted by B_n . Note that from $\partial_{n-1} \circ \partial_n = 0$ it follows that $B_n \subseteq Z_n$ for all n .

The n -th simplicial homology group of X , with coefficients in k , is the vector space $H_n := Z_n / B_n$.

Persistent homology is the homology of a *filtration*, i.e. an increasing sequence of simplicial complexes

$$X_0 \subset X_1 \subset \dots \subset X_n = X,$$

as opposed to that of a single simplicial complex.

It assigns to a filtration the homology groups of the simplicial complexes $H_n(X_v)$ and the linear maps $i_{v,w} : H_n(X_v) \rightarrow H_n(X_w)$ induced in homology by the inclusions $X_v \rightarrow X_w$ for all $v \leq w$. Note that the linear maps $i_{v,v+1}$ are not always injective, meaning that some homological features can disappear along the filtration. These features are encoded by the persistent homology generators: an element $g \in H_n(X_v)$ such that there is no $h \in H_n(X_w)$ for $w < v$ with the property that $i_{w,v} h = g$. Two indices completely determine a generator $g \in H_n(X)$, namely its birth, β_g and its death δ_g . The index β_g traces the first index such that g is in the filtration and δ_g is the index of the simplicial complex in which the cycle becomes a boundary (i.e. disappears homologically). The persistence (lifetime) of a generator is measured by $p_g := \delta_g - \beta_g$. The length of a cycle, that is the number of faces composing it, is denoted by λ_g .

For each homology group, the information about the filtration is collected in a barcode: the set of intervals $[\beta_g; \delta_g]$ for all generators

$g \in H_n$, which constitutes a handy complete invariant of H_n [16]. An alternative way to represent the persistent homology of a filtration is through persistence diagrams [16,50], which we use extensively in the SI. A persistence diagram is a set of points in the plane counted with multiplicity. It can be recovered from the barcode considering the points $(\beta_g, \delta_g) \in \mathbb{R}^2$ with multiplicity given by the number of generators with the same persistence interval. In the SI, the reader can find H_1 persistent diagrams of the real world datasets examined for the classification, together with the explicit comparison to the results for their relevant randomized versions.

Filtrations

In classical applications, the filtration is obtained from a point cloud using the Rips-Vietoris complex and persistent homology used to uncover robust topological features of the point cloud. We instead use the clique weight rank filtration to uncover properties deriving from the topology and weighted structure of weighted networks.

Recalling that an n -clique is a complete subgraph on $n+1$ vertices, the *clique complex* is a simplicial complex built from the cliques of a graph. Namely there is a n -face in the simplicial complex for every $(n+1)$ -clique in the graph. The compatibility relations are satisfied because subsets of cliques and intersection of cliques are cliques themselves.

The *Weight Rank Clique filtration* on a weighted network Ω combines the clique complex construction with a thresholding on weights following three main steps.

- Rank the weights of links from ω_{max} to ω_{min} : the discrete parameter ϵ_t indexes the sequence.
- At each step t of the decreasing edge ranking we consider the thresholded graph $G(\omega_{ij}, \epsilon_t)$, i.e. the subgraph of Ω with links of weight larger than ϵ_t .
- For each graph $G(\omega_{ij}, \epsilon_t)$ we build the clique complex $K(G, \epsilon_t)$.

The clique complexes are nested along the growth of t and determine the weight rank clique filtration. Note that this construction is in fact the clique complex of each element in the graph filtration.

In particular, persistent one dimensional cycles in the weight rank clique filtration represent weighted loops with much weaker internal links.

There is a conceptual difference in interpreting H_1 persistent homology of data with the Rips-Vietoris filtration and H_1 persistent homology of weighted networks with the weight rank clique filtration. While in the first case persistent generators are relevant and considered features of the data, short cycles are more interesting for networks. This is because random networks, or randomisations of real networks, display one dimensional persistent generators at all scales, while short lived generators testify the presence of local organisation properties on different scales.

Computational Complexity

Computing the filtration of a large dataset can be extremely demanding computationally. The identification of the maximal cliques requires in general exponential time, although algorithms exists for special cases that allow solutions to be obtained in polynomial time. In addition, the javaPlex library [51] requires the explicit enumeration of the simplicial facets appearing at each filtration step, which implies the need for large memory resources in order to calculate the persistent homology. However, there are a number of simplifications and improvements to the brute force approach that provide a significant reduction of the problem's complexity. In the metrical case, this is usually done by constructing a smaller complex, the *witness complex* [52],

which approximates with controlled precision [52] the homology of the original data.

In the case of non-metrical discrete spaces, for example networks, one cannot easily construct a witness complex through a controlled sub-sampling of the network. Luckily, it is still possible to reduce the computational complexity in different ways: first, one can limit the analysis to the first s homology groups, which amounts to restricting the clique detection and storage to cliques up to size $s+2$, which reduces the problem to polynomial in time and memory; second, it is possible to parallelize the computation of persistent homology [53]; finally, the more elegant solution is to calculate the homology of an homologically equivalent but much smaller filtration (see the tidy set construction [54]). With respect to the standard clique complex case, the tidy set in particular was shown to reduce the number of simplices along the filtration of various orders of magnitude number of simplices and of one order of magnitude the total memory required. Therefore, a combina-

tion of the techniques mentioned above allows to scale up dataset sizes to large-scale networks.

Supporting Information

File S1
(PDF)

Acknowledgments

The authors acknowledge M. Rasetti for stimulating discussions.

Author Contributions

Conceived and designed the experiments: GP FV. Performed the experiments: GP MS ID FV. Analyzed the data: GP MS ID FV. Contributed reagents/materials/analysis tools: GP MS ID FV. Wrote the paper: GP MS ID FV.

References

- Newman MEJ (2003) The structure and function of complex networks. *SIAM Rev* 45: 167–256.
- Boccaletti S, Latora V, Moreno Y, Chavez M, Hwang DH (2006) Complex networks: Structure and dynamics. *Phys Rep* 424: 175–308.
- Dorogovtsev SN, Goltsev AV, Mendes JFF (2008) Critical phenomena in complex networks. *Rev Mod Phys* : 1275–1335.
- Barrat A, Barthélemy M, Pastor-Satorras R, Vespignani A (2004) The architecture of complex weighted networks. *Proc Natl Acad Sci USA* 101: 3747–3752.
- Barabási A (1999) Emergence of scaling in random networks. *Science* 286: 509–512.
- Milo R, Shen-Orr S, Itzkovitz S, Kashtan N, Chklovskii D, et al. (2002) Network motifs: Simple building blocks of complex networks. *Science* 298: 824–827.
- Vázquez A, Dobrin R, Sergi D, Eckmann JP, Olvtai ZN, et al. *Proc Natl Acad Sci USA* 101: 17940–17945.
- Mahadevan P, Krioukov D, Fall K, Vahdat A (2006) Systematic topology analysis and generation using degree correlations. *ACM SIGCOMM* 36: 135–146.
- Conradi C, Flockerzi D, Raisch J, Stelling J (2007) Subnetwork analysis reveals dynamic features of complex (bio)chemical networks 104: 19175–19180.
- Egúiluz VM, Chialvo DR, Cecchi GA, Baliki M, Apkarian AV (2005) Scale-free brain functional networks. *Phys Rev Lett* 92: 028102.
- Song WM, Di Matteo T, Aste T (2012) Hierarchical information clustering by means of topologi3 cally embedded graphs. *PLoS One* 7, e31929.
- Tumminello M, Aste T, Di Matteo T, Mantegna RN (2005) A tool for filtering information in complex systems. *Proc Natl Acad Sci USA* 102: 10421–10426.
- Chalupa J, Leath PL, Reich GR (1979) Bootstrap percolation on a bethe lattice. *J Phys C*.
- Serrano M, Boguñá M, Vespignani A (2009) Extracting the multiscale backbone of complex weighted networks. *Proc Natl Acad Sci USA* 106: 6483.
- Ghrist R (2008) Barcodes: The persistent topology of data. *B AM Math Soc* 45.
- Carlsson G, Zomorodian A (2005) Persistent homology - a survey. *Discrete Comput Geom* 33: 249–274.
- Carlsson G (2009) Topology and data. *B Am Math Soc* 46: 255–308.
- Petri G, Scolamiero M, Donato I, Vaccarino F (2013) Metric and weighted clique persistent ho3 mology for complex networks. In: *Proceedings of the European Conference on Complex Systems 2012*.
- Lum P, Singh G, Lehman A, Ishkanov T, Vejdemo-Johansson M, et al. *Scientific Reports*.
- Hyeckyoung L, Chung M, Hyejin K, Bung-Nyun K, L DS (2011) Discriminative persistent homology of brain networks. In: *Biomedical Imaging: From Nano to Macro, 2011 IEEE International Symposium on*. 841–844. doi:10.1109/ISBL.2011.5872535.
- Nicolau M, Levine A, Carlsson G (2011) Topology based data analysis identifies a subgroup of breast cancers with a unique mutational profile and excellent survival. *Proceedings of the National Academy of Sciences* 108: 7265–7270.
- Weygaert R, Vegter G, Edelsbrunner H, Jones B, Pranav P, et al. (2011) Alpha, betti and the megaparsec universe: On the topology of the cosmic web. In: Gavrilova M, Tan C, Mostafavi M, editors, *Transactions on Computational Science XIV*, Springer Berlin Heidelberg, volume 6970 of *Lecture Notes in Computer Science*. 60–101. doi:10.1007/978-3-642-25249-5_3. URL http://dx.doi.org/10.1007/978-3-642-25249-5_3.
- Opsahl T, Colizza V, Panzarasa P, Ramasco JJ (2008) Prominence and control: The weighted rich-club effect. *Phys Rev Lett* 101: 168702.
- Pajevic D, Plenz S (2012) The organization of strong links in complex networks. *Nat Phys* 8: 429–436.
- Barthélemy M (2011) Spatial networks. *Phys Rep* 499: 1–101.
- Antonioni A, Tomassini M (2012) Degree correlations in random geometric graphs. *Phys Rev E* 86: 037101.
- Barrat A, Barthélemy M, Vespignani A (2005) The effects of spatial constraints on the evolution of weighted complex networks. *J Stat Mech* 05: P05003.
- Palla G, Derényi I, Farkas I, Vicsek T (2005) The effects of spatial constraints on the evolution of weighted complex networks.
- Stehlé J, Voirin N, Barrat M, Cattuto C, Isella L, et al. (2011) High-resolution measurements of face-to-face contact patterns in a primary school. *PLoS One* 6: e23176.
- Gfeller D, De Los Rios P (2007) Spectral coarse graining of complex networks. *Phys Rev Lett* 99: 38701.
- Chavez M, Hwang DU, Amann A, Hentschel HGE, Boccaletti S (2005) Synchronization is enhanced in weighted complex networks. *Phys Rev Lett* 94: 218701.
- Farkas IJ, Derényi I, Barabási AL, Vicsek T (2001) Spectra of ‘real-world’ graphs: Beyond the semicircle law. *Phys Rev E* 64.
- Jun WU, Barahona M, Yue-Jin T, Hong-Zhong D (2010) Natural connectivity of complex networks. *Chin Phys Lett* 27: 078902.
- Barahona M, Pecora LM (2002) Synchronization in small-world systems. *Phys Rev Lett* 89: 054101.
- MacArthur BD, Sánchez-García RJ (2009) Spectral characteristics of network redundancy. *Phys Rev E* 80: 026117.
- Colizza V, Flammini A, Serrano MA, Vespignani A (2006) Detecting rich-club ordering in complex networks. *Nat Phys* 2: 110–115.
- Boguñá M PF, Krioukov D (2010) Sustaining the internet with hyperbolic mapping. *Nat Comms* 1: 1–8.
- Grady D, Thiemann C, Brockmann D (2012) Robust classification of salient links in complex networks. *Nat Comm* 3: 864.
- Carlsson G, Zomorodian A (2009) Theory of multidimensional persistence. *Discr Comput Geom* 42: 71–93.
- Watts D, Strogatz SH (1998) Collective dynamics of ‘small-world’ networks. *Nature* 393: 440–442.
- Opsahl T, Panzarasa P (2009) Clustering in weighted networks. *Soc Net* 31: 155–163.
- Opsahl T (2010) Triadic closure in two-mode networks: Redefining the global and local clustering coefficients. *Soc Net*.
- Gardner T, di Bernardo D, Lorenz D, Collins J (2003) Inferring genetic networks and identifying compound mode of action via expression profiling. *Science* 301: 102–105.
- Newman MEJ (2001) The structure of scientific collaboration networks. *Proc Natl Acad Sci USA* 98: 404–409.
- Petri G (2013). Holes - python package for persistent homology calculations. URL <http://lordgrilo.github.com/Holes/>.
- Barthélemy M (2011) Spatial networks. *Phys Rep* 499.
- Penrose M (2003) *Random Geometric Graphs*. Oxford, UK: Oxford University Press.
- Palla Gea (2007) Directed network modules. *New J Phys* 9.
- Munkres JR (1984) *Elements of Algebraic Topology*. 2725 Sand Hill Road Menlo Park, California 94025: Addison-Wesley Publishing Company.
- Steiner DC, Edelsbrunner H, Harer J (2007) Stability of persistence diagrams. *Discrete Comput Geom* 37: 103–120.
- Tausz A, Vejdemo-Johansson M, Adams H (2011). Javaplex: A research software package for persistent (co)homology. “Software available at <http://code.google.com/javaplex/>”.
- de Silva V, Carlsson G (2004) Topological estimation using witness complexes. *Symp Point- Based Graphics, ETH Zurich*.
- Bauer U, Kerber M, Reininghaus J (2013) Clear and compress: Computing persistent homology in chunks. *arXiv: 13030477*.
- Zomorodian A (2010) The tidy set: a minimal simplicial set for computing homology of clique complexes. *Proceedings of the 2010 annual symposium on Computational geometry* : 257–266.